Semi-Deterministic Broadcast Channels with Cooperation

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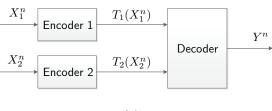
December, 2014

Outline

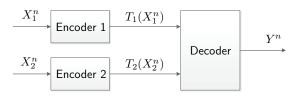
- Motivation and past work
- AK problem with one-sided encoder cooperation
- SD-BC with one-sided decoder cooperation
- Duality
- Summary

• The two-encoder multiterminal source coding problem [Berger, 1978], [Tung, 1978].

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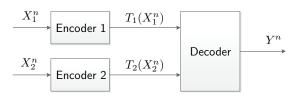
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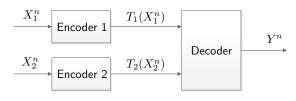
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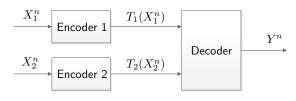
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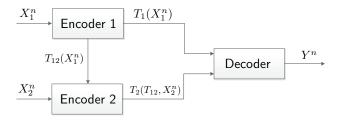
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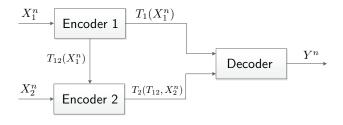


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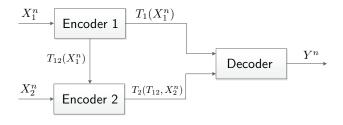
- Ahlswede-Körner (AK) problem (1975).
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 - Milestone towards multiuser channel-source duality.



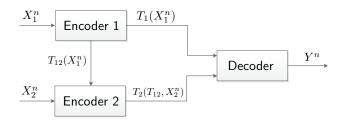
Without cooperation [Ahlswede-Körner, 1975]



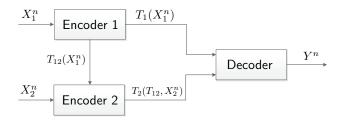
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- Decoder Output: $(X_1^n, X_2^n, Y^n) \in \mathcal{T}_{\epsilon}^{(n)} (P_{X_2} P_{Y|X_2} \mathbb{1}_{\{X_1 = f(Y)\}}).$

Theorem (Coordination-Capacity Region)

For a desired coordination distribution $P_{X_2}P_{Y|X_2}\mathbb{1}_{\{X_1=f(Y)\}}$:

$$\mathcal{C}_{AK} = \bigcup \left\{ \begin{array}{c} R_{12} \geq I(V; X_1) - I(V; X_2) \\ R_1 \geq H(X_1 | V, U) \\ R_2 \geq I(U; X_2 | V) - I(U; X_1 | V) \\ R_1 + R_2 \geq H(X_1 | V, U) + I(V, U; X_1, X_2) \end{array} \right\}$$

where the union is over all $P_{X_1,X_2}P_{V|X_1}P_{U|X_2,V}P_{Y|X_1,U,V}$ with $P_{X_2}P_{Y|X_2}\mathbb{1}_{\{X_1=f(Y)\}}$ as marginal.

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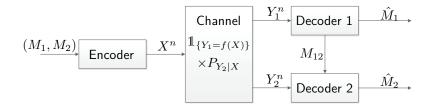
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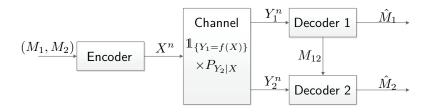
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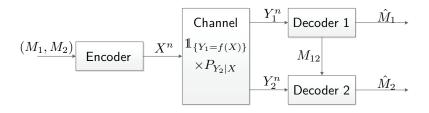
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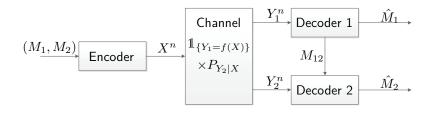
Without cooperation [Gelfand and Pinsker, 1980]



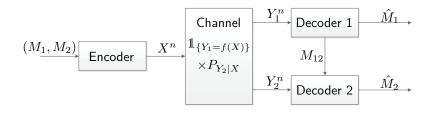
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- Decoders' Output: $\hat{M}_1(Y_1^n)$ and $\hat{M}_2(M_{12},Y_2^n)$.

Semi-Deterministic BC with Cooperation - Solution

Theorem (Capacity Region)

The capacity region is:

$$\mathcal{C}_{BC} = \bigcup \left\{ \begin{array}{c} R_{12} \geq I(V;Y_1) - I(V;Y_2) \\ R_1 \leq H(Y_1) \\ R_2 \leq I(V,U;Y_2) + R_{12} \\ R_1 + R_2 \leq H(Y_1|V,U) + I(U;Y_2|V) + I(V;Y_1) \end{array} \right\}$$

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Achievability via rate splitting, Marton coding and Wyner-Ziv-like coding for cooperation protocol.

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Generalization of [Lapidoth and Wang, 2013].

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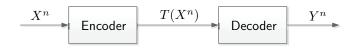
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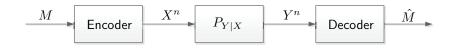
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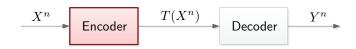
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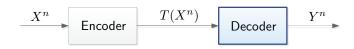
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- ullet Solving one problem \implies Valuable insight into solving dual.



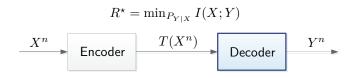


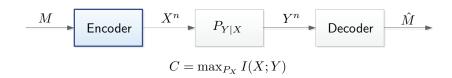




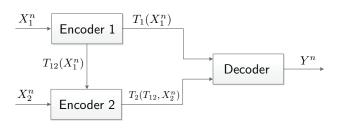


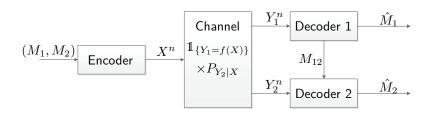




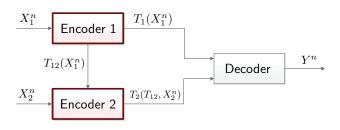


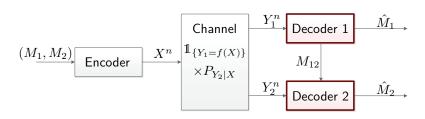
AK Problem vs. Semi-Deterministic BC:



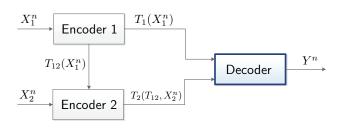


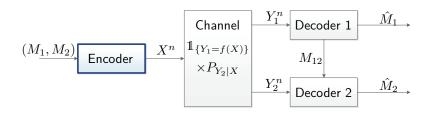
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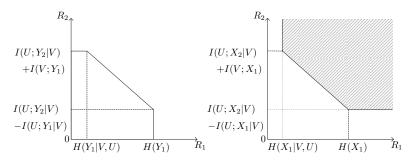
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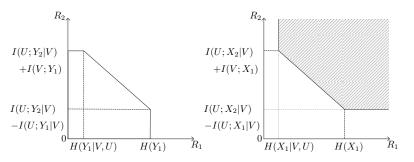
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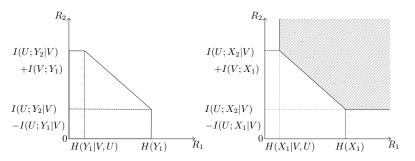
AK Problem

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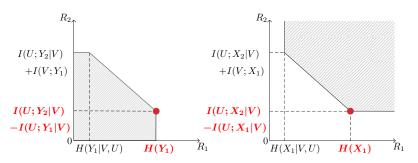




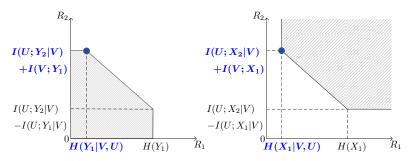
Semi-Deterministic BC with Cooperation	Ahlswede-Körner Problem with Cooperation
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(R_1,R_2) at Lower Corner Point:	(R_1,R_2) at Lower Corner Point:
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$\Big(H(Y_1 V,U)\;,\;I(U;Y_2 V)+I(V;Y_1)\Big)$	$\Big(H(X_1 V,U)\;,\;I(U;X_2 V)+I(V;X_1)\Big)$

• AK problem with cooperation.

- AK problem with cooperation.
- SD-BC with cooperation.

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- SD-BC with cooperation.
- Duality:

- AK problem with cooperation.
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- Duality:
 - Transformation principles.

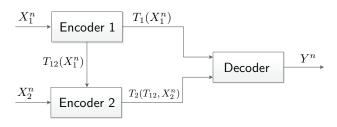
- AK problem with cooperation.
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- Duality:
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 - Corner point correspondence.

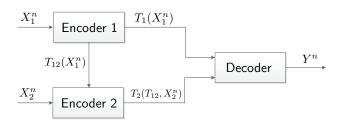
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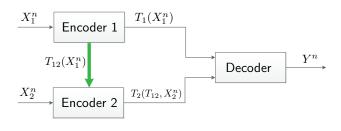
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Thank you!



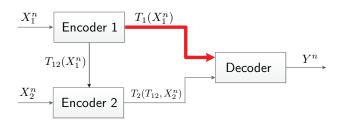


Rate	Corner Point 1	Corner Point 2
R_{12}	$I(V;X_1) - I(V;X_2)$	$I(V;X_1) - I(V;X_2)$
R_1	$H(X_1)$	$H(X_1 V,U)$
R_2	$I(U; X_2 V) - I(U; X_1 V)$	$I(U; X_2 V) + I(V; X_1)$



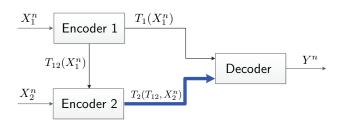
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• Cooperation: Wyner-Ziv scheme to convey V^n via cooperation link.



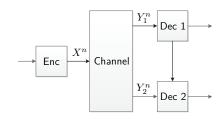
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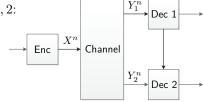


Rate	Corner Point 1	Corner Point 2
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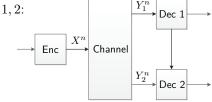
- Cooperation: Wyner-Ziv scheme to convey V^n via cooperation link.
- Corner Point 1: V^n is transmitted to dec. by Enc. 1 within X_1^n .
- ullet Corner Point 2: V^n is explicitly transmitted to dec. by Enc. 2.



- Rate Splitting: $M_i = (M_{i0}, M_{ij}), j = 1, 2$:
 - (M_{10}, M_{20}) Public message;



- Rate Splitting: $M_j = (M_{j0}, M_{jj}), j = 1, 2$:
 - (M_{10}, M_{20}) Public message;
 - (M_{11}, M_{22}) Private messages.

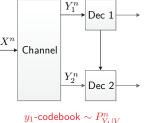


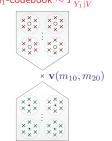
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• Codebook Structure: Marton (with common message).

- Rate Splitting: $M_j = (M_{j0}, M_{jj}), j = 1, 2$:
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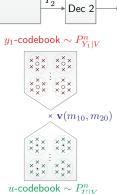
u-codebook $\sim P_{U|V}^n$

Enc

- Rate Splitting: $M_j = (M_{j0}, M_{jj}), j = 1, 2$:
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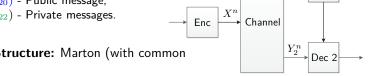
• Codebook Structure: Marton (with common message). Channel Y_2^n Dec 2

Cooperation:

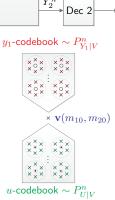


Dec 1

- Rate Splitting: $M_i = (M_{i0}, M_{ij}), j = 1, 2$:
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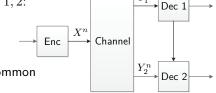


- Codebook Structure: Marton (with common message).
- Cooperation:
 - 1. Partition common message c.b into $2^{nR_{12}}$ bins.



Dec 1

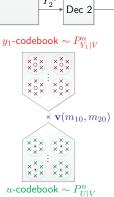
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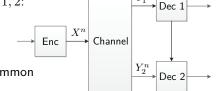
 Codebook Structure: Marton (with common message).

Cooperation:

- 1. Partition common message c.b into $2^{nR_{12}}$ bins.
- 2. Convey bin number via link.



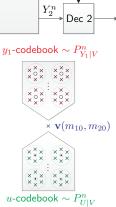
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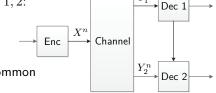
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Cooperation:

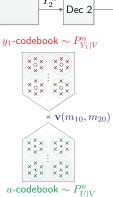
- 1. Partition common message c.b into $2^{nR_{12}}$ bins.
- 2. Convey bin number via link.
- Gain at Dec. 2:



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- Codebook Structure: Marton (with common message).
- Cooperation:
 - 1. Partition common message c.b into $2^{nR_{12}}$ bins.
 - 2. Convey bin number via link.
- Gain at Dec. 2: Reduced search space of common message c.w by R_{12} .



Via telescoping identities:

1. Auxiliaries: $V_i = (M_{12}, Y_1^{i-1}, Y_{2,i+1}^n)$ and $U_i = M_2$.

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$$= \sum_{i=1}^{n} \left[I(M_{2}; Y_{2,i}^{n} | M_{12}, Y_{1}^{i-1}) - I(M_{2}; Y_{2,i+1}^{n} | M_{12}, Y_{1}^{i}) \right] + I(M_{2}; M_{12})$$

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Via telescoping identities:

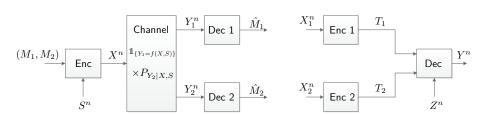
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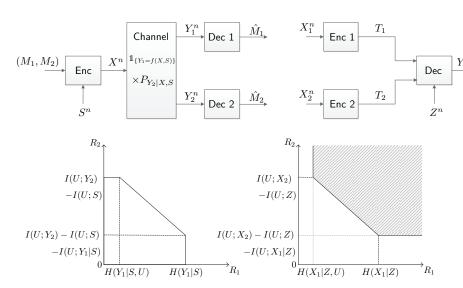
Replaces 4 uses of Csiszár Sum Identity!

State-Dependant Semi-Deterministic BC vs. Dual:

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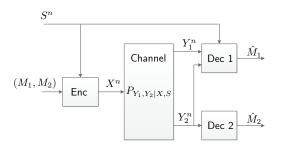


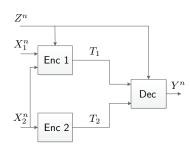
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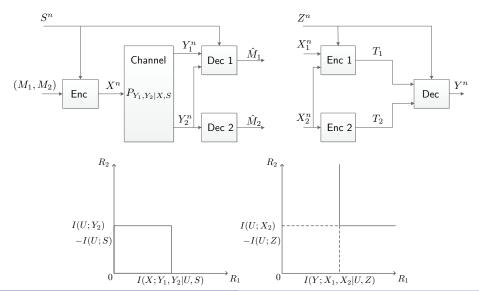
State-Dependant Output-Degraded BC vs. Dual:

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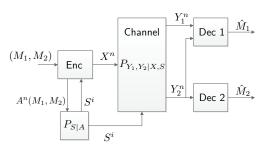


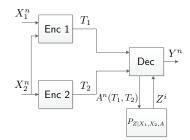
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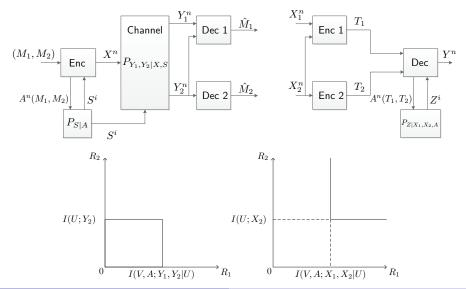
Action-Dependant Output-Degraded BC vs. Dual:

Action-Dependant Output-Degraded BC vs. Dual:





Action-Dependant Output-Degraded BC vs. Dual:



Achieving Corner Point 1:

$$(I(V; X_1|X_2), H(X_1), I(U; X_2|X_1, V)).$$

Achieving Corner Point 1:

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- Encoder 1 to Decoder: Conveys X_1^n to the decoder in a lossless manner.
- Encoder 2 to Decoder: The decoder knows X_1^n and therefore V^n . Wyner-Ziv coding to convey U^n .

Achieving Corner Point 2:

$$(I(V; X_1|X_2), H(X_1|V,U), I(U; X_2|V) + I(V; X_1)).$$

Achieving Corner Point 2:

$$(I(V; X_1|X_2), H(X_1|V,U), I(U; X_2|V) + I(V; X_1)).$$

Cooperation: Same.

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- Cooperation: Same.
- Encoder 2 to Decoder: Knows V^n . Conveys the index of V^n and uses superposition coding to convey U^n .

Achieving Corner Point 2:

$$(I(V; X_1|X_2), \frac{H(X_1|V,U)}{I(U;X_2|V)}, I(U;X_2|V) + I(V;X_1)).$$

- Cooperation: Same.
- Encoder 2 to Decoder: Knows V^n . Conveys the index of V^n and uses superposition coding to convey U^n .
- Encoder 1 to Decoder: The decoder knows (V^n, U^n) . Binning scheme to convey X_1^n in a lossless manner.

AK Problem with Cooperation - Proof Outline

Converse:

AK Problem with Cooperation - Proof Outline

Converse:

Standard techniques while defining

$$V_i = (T_{12}, X_1^{n \setminus i}, X_{2,i+1}^n),$$

 $U_i = T_2,$

for every $1 \le i \le n$.

AK Problem with Cooperation - Proof Outline

Converse:

Standard techniques while defining

$$V_i = (T_{12}, X_1^{n \setminus i}, X_{2,i+1}^n),$$

 $U_i = T_2,$

for every $1 \le i \le n$.

Time mixing properties.

• Rate Splitting: $M_j = (M_{j0}, M_{jj}), \ j = 1, 2:$ $\longrightarrow \boxed{\text{Enc}} \xrightarrow{X^n} \text{Channel} \xrightarrow{Y_1^n} \boxed{\text{Dec 1}}$

• Rate Splitting: $M_j = (M_{j0}, M_{jj})$, j = 1, 2:
• (M_{10}, M_{20}) - Public message;

Enc X^n Channel Y_2^n Dec 1

• Rate Splitting: $M_j = (M_{j0}, M_{jj})$, j = 1, 2:

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• Enc X^n Channel Y_2^n Dec 2

- $\xrightarrow{Y_1^n}$ Dec 1 • Rate Splitting: $M_i = (M_{i0}, M_{ii}), j = 1, 2$: • (M_{10}, M_{20}) - Public message; • (M_{11}, M_{22}) - Private messages. → Enc • Codebook Structure: Marton:

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• Channel

• Codebook Structure: Marton:

Goldfeld/Permuter/Kramer

▶ Public Message: $(M_{10}, M_{20}) \longrightarrow V^n$.

- Rate Splitting: $M_i = (M_{i0}, M_{ij}), j = 1, 2$: $\xrightarrow{Y_1^n}$ Dec 1 \blacktriangleright (M_{10}, M_{20}) - Public message; • (M_{11}, M_{22}) - Private messages. → Enc • Codebook Structure: Marton:
 - - ▶ Public Message: $(M_{10}, M_{20}) \longrightarrow V^n$.
 - ▶ Private Messages Superposed on V^n :

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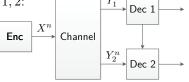
 (M_{11}, M_{22}) Private messages.

 Channel

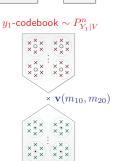
 Codebook Structure: Marton:
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 - 1. $M_{11} \longrightarrow Y_1^n$;

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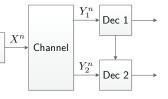
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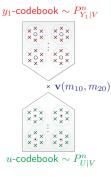
u-codebook $\sim P_{U|V}^n$

→ Enc

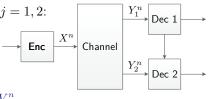
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- Codebook Structure: Marton:
 - ▶ Public Message: $(M_{10}, M_{20}) \longrightarrow V^n$.
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 - 1. $M_{11} \longrightarrow Y_1^n$;
 - 2. $M_{22} \longrightarrow U^n$.
- Decoding: Joint typicality decoding.

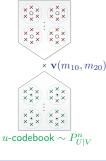


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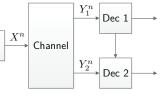
 y_1 -codebook $\sim P_{Y_1|V}^n$

- Codebook Structure: Marton:
 - ▶ Public Message: $(M_{10}, M_{20}) \longrightarrow V^n$.
 - ▶ Private Messages Superposed on V^n :
 - 1. $M_{11} \longrightarrow Y_1^n$;
 - 2. $M_{22} \longrightarrow U^n$.
- Decoding: Joint typicality decoding.
- Cooperation: Bin number of V^n $2^{nR_{12}}$ bins.

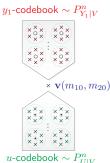


→ Enc

- Rate Splitting: $M_j = (M_{j0}, M_{jj}), j = 1, 2$:
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 - (M_{11}, M_{22}) Private messages.



- Codebook Structure: Marton:
 - ▶ Public Message: $(M_{10}, M_{20}) \longrightarrow V^n$.
 - ▶ Private Messages Superposed on V^n :
 - 1. $M_{11} \longrightarrow Y_1^n$;
 - 2. $M_{22} \longrightarrow U^n$.
- Decoding: Joint typicality decoding.
- Cooperation: Bin number of V^n $2^{nR_{12}}$ bins.
- Gain: Dec. 2 reduces search space of V^n by R_{12} .

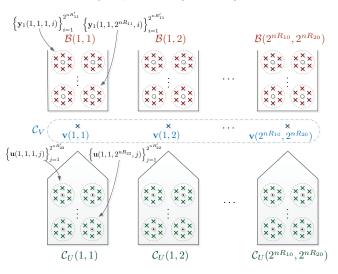


Semi-Deterministic BC with Cooperation - Proof Outline

Achievability: Split $M_i = (M_{i0}, M_{ii})$, i = 1, 2. Code construction:

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Legend:

- Private message m_{11}
- Private message m_{22}
- \times v-codeword ($\sim P_V$) \times - y_1 -codeword ($\sim P_{Y_1}$)
- \times u-codeword ($\sim P_{U|V}$)