# MIMO Gaussian Broadcast Channels with Common, Private and Confidential Messages 

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Ben Gurion University
IEEE Information Theory Workshop
September, 2016

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## Motivation - Banking Site



| Log in Onine Banking |
| :--- |
| BHI Online |
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| New York |
| Israe! |
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| Security and Privacy |

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- Investor Relations
- Sustainability and Social Responsibility
- Reports and Forecasts
- Awards and Recogrition


Bank Hapoalim
Your Gateway to Israel

## Bank Hapoalim Announces Second Quarter 2016

Net Profit totaled NIS 1,117 million, Return on Equity of $13.9 \%$ Cash Dividend Payout of NIS 223 million

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| Indices |  |  |
| :---: | :---: | :---: |
| - Dow | 18419.3 | 0.10\% |
| - Nasdaq | 5227.21 | 0.27\% |
| - SP500 | 2170.88 | 0.00\% |
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| * Quotes delayed by at least 15 minutes |  |  |
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## Common

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- Security Criterion:

$$
\frac{1}{n} I\left(M_{1} ; \mathbf{Y}_{2}^{n}\right) \xrightarrow[n \rightarrow \infty]{ } 0
$$

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Q: Do Gaussian inputs achieve boundary points?

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* Solution for two last unsolved cases via Upper Concave Envelopes $\star$


## MIMO Gaussian BC - Secrecy-Capacity Results

Without a Common Message: $M_{1}$ - Confidential ; $M_{2}$ - Private

## Theorem (ZG 2016)

The secrecy-capacity region for a covariance constraint $\mathrm{K} \succeq 0$ is

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\hat{\mathcal{C}}_{\mathrm{K}}=\bigcup_{0 \preceq \mathrm{~K}^{\star} \preceq \mathrm{K}}\left\{\left(R_{1}, R_{2}\right) \in \mathbb{R}_{+}^{2} \left\lvert\, \begin{array}{l}
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- $\boldsymbol{R}_{\mathbf{2}}$ Bound - Capacity of MIMO Gaussian PTP to User 2: Input covariance $\mathrm{K}-\mathrm{K}^{\star}$; Noise covariance $\mathrm{I}+\mathrm{G}_{2} \mathrm{~K}^{\star} \mathrm{G}_{2}^{\top}$.


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- $R_{0}$ Bound: MIMO Gaussian PTP with remaining covariance $K-\left(K_{1}+K_{2}\right)\left(K_{1}, K_{2}\right.$ are noises $)$.


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## Thank You!

