

# Maximum Likelihood Approach for Symmetric Distribution Property Estimation

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# Property estimation

- $p$ : unknown discrete distribution over  $k$  elements
- $f(p)$ : a property of  $p$
- $\epsilon$ : accuracy parameter,  $\delta$ : error probability
- Given: access to independent samples from  $p$
- Goal: estimate  $f(p)$  to  $\pm\epsilon$  with probability  $> 1 - \delta$

Usually,  $\delta$  constant, say **0.1**

Focus of the talk: **large  $k$**  (really large)

# Sample complexity

$X_1, X_2, \dots, X_n$ : independent samples from  $p$

$\hat{f}(X_1^n)$ : estimate of  $f(p)$

Sample complexity of  $\hat{f}(X_1^n)$ :

$$S(\hat{f}, \varepsilon, k, \delta) = \min\{n: \forall p, \Pr(|\hat{f}(X_1^n) - f(p)| \geq \varepsilon) \leq \delta\}$$

Sample complexity of  $f$ :

$$S(f, \varepsilon, k, \delta) = \min_{\hat{f}} S(\hat{f}, \varepsilon, k, \delta)$$

# Symmetric properties

Permuting labels does not change  $f(p)$

Examples:

- $H(p) \triangleq -\sum_x p_x \cdot \log(p_x)$
- $S(p) \triangleq \sum_x 1_{p_x > 0}$

Renyi entropy, distance to uniformity, unseen symbols, divergences, etc

# Sequence maximum likelihood

$N_x$ : # times  $x$  appears in  $X_1^n$  (multiplicity)

$p^e$ : empirical distribution

$$p_x^e = \frac{N_x}{n}$$

$$f^e(X_1^n) = f(p^e)$$

Empirical estimators are maximum likelihood estimators

$$p^e = \arg \max_p p(X_1^n)$$

Call this sequence maximum likelihood **(SML)**

# Empirical entropy estimation

Empirical estimator:

$$H^e(X_1^n) = H(p^e) = \sum_x \frac{N_x}{n} \log \frac{n}{N_x}.$$
$$S(H^e, \varepsilon, k, 0.1) = \Theta\left(\frac{k}{\varepsilon}\right)$$

Various corrections proposed:

Miller-Maddow, Jackknifed estimator, Coverage adjusted, ...

**Sample complexity:  $\Omega(k)$**

[Paninski'03]:  $S(H, \varepsilon, k, 0.1) = o(k)$  (existential)

**Note:** For  $\varepsilon \ll 1/k$ , empirical estimators are the best

# Entropy estimation

[ValiantValiant'11a]: Constructive proofs based on LP:

$$S(H, \varepsilon, k, 0.1) = \Theta_{\varepsilon} \left( \frac{k}{\log k} \right)$$

- [YuWang'14, HanJiaoVenkatWeissman'14, ValiantValiant11b]:

Simplified algorithms, and exact rates:

$$S(H, \varepsilon, k, 0.1) = \Theta \left( \frac{k}{\varepsilon \log k} \right)$$

# Support coverage

Expected number of symbols when  $p$  is sampled  $m$  times

$$S_m(p) = \sum_x (1 - (1 - p_x)^m)$$

Goal: Estimate  $S_m(p)$  to  $\pm(\varepsilon \cdot m)$

[OrlitskySureshWu'16, ZouValiantetal'16]:

$$n = \Theta\left(\frac{m}{\log m} \log\left(\frac{1}{\varepsilon}\right)\right) \text{ samples for } \delta = 0.1$$

# Known results summary

Many symmetric properties: entropy, support size, distance to uniform, support coverage

- Different estimator for each property
- Sophisticated results from approximation theory

# Main result

Simple, ML based plug-in method that is sample-optimal for entropy, support-coverage, distance to uniform, support size.

Property	Notation	SML	Optimal	References	PML
Entropy	$H(p)$	$\frac{k}{\varepsilon}$	$\frac{k}{\log k} \frac{1}{\varepsilon}$	[VV11a, WY16, JVHW15]	optimal <sup>1</sup>
Support size	$\frac{S(p)}{k}$	$k \log \frac{1}{\varepsilon}$	$\frac{k}{\log k} \log^2 \frac{1}{\varepsilon}$	[WY15]	optimal
Support coverage	$\frac{S_m(p)}{m}$	$m$	$\frac{m}{\log m} \log \frac{1}{\varepsilon}$	[OSW16]	optimal
Distance to $u$	$\ p - u\ _1$	$\frac{k}{\varepsilon^2}$	$\frac{k}{\log k} \frac{1}{\varepsilon^2}$	[VV11b, JHW16]	optimal

# Profiles

**Profile** of a sequence is the **multiset** of **multiplicities**:

$$\Phi(X_1^n) = \{N_x\}$$

$$X_1^n = (1H, 2T), \text{ or } X_1^n = (2H, 1T), \Phi(X_1^n) = \{1,2\}$$

Symmetric properties depend on **multiset** of probabilities

Coins w/ bias 0.4, and 0.6 have same symmetric property

**Optimal estimators** have **same** output for sequences with **same profile**.

**Profiles are sufficient statistic**

# Profile maximum likelihood [OSVZ'04]

Probability of a profile:

$$p(\Phi(X_1^n)) = \sum_{Y_1^n: \Phi(Y_1^n) = \Phi(X_1^n)} p(Y_1^n)$$

**Maximize the profile probability:**

$$p_{\Phi}^{PML} = \arg \max_p p(\Phi(X_1^n))$$

$X_1^n = (1H, 2T)$ :

SML:  $(2/3, 1/3)$

PML:  $(1/2, 1/2)$

# PML for symmetric properties

To estimate a symmetric  $f(p)$ :

- Find  $p^{PML}(\Phi(X_1^n))$
- Output  $f(p^{PML})$

Advantages:

- No tuning parameters
- Not function specific

# Main result

PML is **sample-optimal** for entropy, support coverage, distance to uniformity, and support size.

# Ingredients

## Guarantee for PML.

If  $n = S(f, \varepsilon, k, \delta)$ , then  $S\left(f(p^{PML}), 2\varepsilon, k, \delta \cdot \frac{e^{3\sqrt{n}}}{10}\right) \leq n$

If  $n = S(f, \varepsilon, k, e^{-3\sqrt{n}})$ , then  $S(f(p^{PML}), 2\varepsilon, k, 0.1) \leq n$

# profiles of length  $n < \frac{e^{3\sqrt{n}}}{10}$

# Ingredients

$n = S(f, \varepsilon, \delta)$ , achieved by an estimator  $\hat{f}(\Phi(X^n))$ :

$p$ : underlying distribution.

- Profiles  $\Phi(X^n)$  such that  $p(\Phi(X^n)) > \delta$ ,

$$p^{PML}(\Phi) \geq p(\Phi) > \delta$$

$$p_{\Phi}^{PML}(\Phi) \geq p(\Phi) > \delta$$

$$|f(p_{\Phi}^{PML}) - f(p)| \leq |f(p_{\Phi}^{PML}) - \hat{f}(\Phi)| + |\hat{f}(\Phi) - f(p)| < 2\varepsilon$$

- Profiles with  $p(\Phi(X^n)) < \delta$ ,

$$p(p(\Phi(X^n)) < \delta) < \delta \cdot \#profiles\ of\ length\ n$$

# Ingredients

**Better error probability guarantees.**

Recall:

$$S(H, \varepsilon, k, 0.1) = \Theta\left(\frac{k}{\varepsilon \cdot \log k}\right)$$

Stronger error guarantees using McDiarmid's inequality:

$$S(H, \varepsilon, k, e^{-n^{0.9}}) = \Theta\left(\frac{k}{\varepsilon \cdot \log k}\right)$$

With twice the samples error drops **exponentially**

Similar results for other properties

# Main result

PML is **sample-optimal** for entropy, unseen, distance to uniformity, and support size.

Even **approximate PML** is optimal for these.

# Algorithms

- EM algorithm [[Orlitsky et al](#)]
- Approximate PML via Bethe Permanents [[Vontobel](#)]
- Extensions of Markov Chains [[VatedkaVontobel](#)]

Polynomial time algorithms for approximating PML

# Summary

- Symmetric property estimation
- PML plug-in approach
  - Universal, simple to state
  - Independent of particular properties
- Directions:
  - Efficient algorithms – for approximate PML
  - Relies heavily on existence of other estimators

# In Fisher's words ...

Of course nobody has been able to prove that maximum likelihood estimates are best under all circumstances. Maximum likelihood estimates computed with all the information available may turn out to be inconsistent. Throwing away a substantial part of the information may render them consistent.

R. A. Fisher

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