

# Capacity Region of the Finite State MAC with Cooperative Encoders and Delayed CSI

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## Motivation (delayed state information)

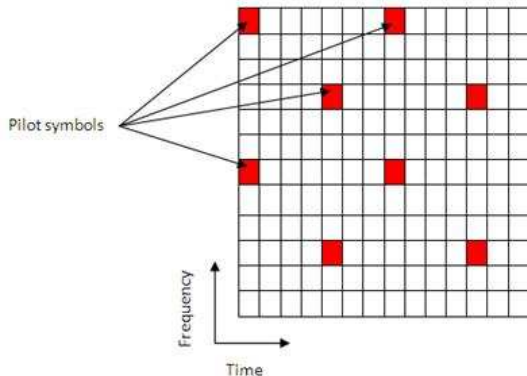
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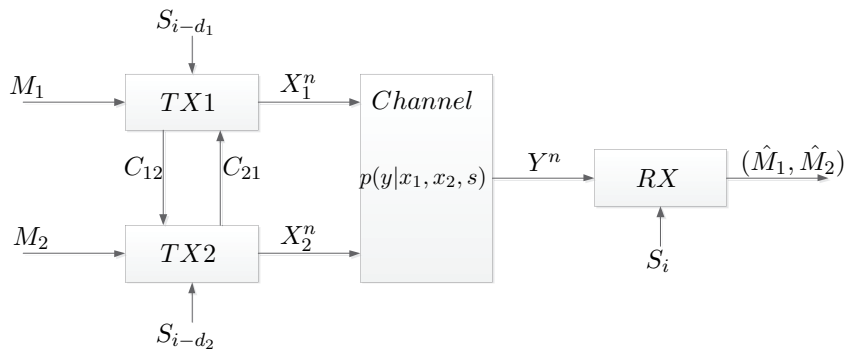
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- Channel state models fading, noise, and interference of uncontrolled signals.
- Channel state information (CSI) needs to be estimated
- In LTE uplink standard, pilot signals are sent by the users

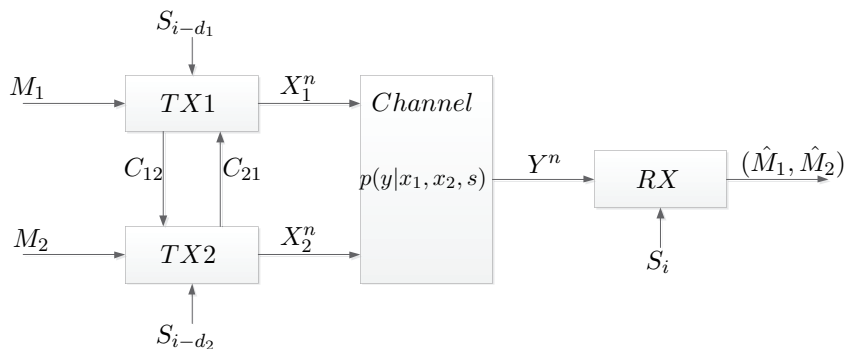


# Uplink with Delayed CSI and conferencing



- CSI known to the Transmitters (TX) and delayed CSI known to the Receivers (RX).
- Conferencing between the TX is possible with limited link.

# Uplink with Delayed CSI and conferencing



- Strictly causal CSI [Steinberg/Lapidoth10][Li/Simeone/Yener10]
- delayed state for Point-to-point case [Viswanathan99]
- No conferencing [Bashar/Shirazi/P 11]
- Assymetrical state [Sen/Alajaji/uksel/Como12]
- Non-causal state at one encoder [Somekh-Baruch/Shamai/Verd'u06] [Kotagiri/Laneman04]

# Channel Model and Notation

- Finite number of states  $\mathcal{S} < \infty$ .
- Channel state is a stationary Markov process independent of the messages.
- The random variables  $S_i, S_{i-d}$  denote the channel state at time  $i$ , and  $i - d$ , respectively.
- The  $(S_i, S_{i-d})$  joint distribution is stationary and is given by

$$P(S_i = s_l, S_{i-d} = s_j) = \pi(s_j)K^d(s_l, s_j).$$

- The channel transition probability at time  $i$  is given by

$$P(y_i | x_{1,i}, x_{2,i}, s_i)$$

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$$P(s^n, v_1^\ell, v_2^\ell) = P(s^n)P(v_1^\ell, v_2^\ell) = \prod_{i=1}^n P(s_i | s_{i-1})P(v_1^\ell, v_2^\ell).$$

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- cooperation link constraint  $C_{12}$  and  $C_{21}$ :

$$\sum_{i=1}^{\ell} \log |\mathcal{V}_{1,i}| \leq nC_{12} ; \quad \sum_{i=1}^{\ell} \log |\mathcal{V}_{2,i}| \leq nC_{21}.$$

## Conferencing encoder

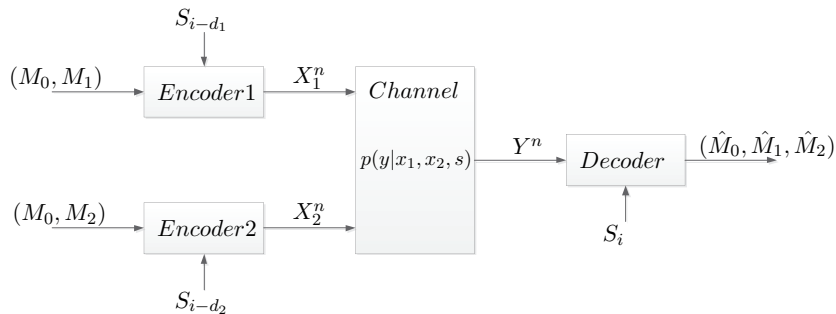
$$\begin{aligned}V_{1,i} &= h_{1,i}(M_1, V_2^{i-1}), \\V_{2,i} &= h_{2,i}(M_2, V_1^{i-1}).\end{aligned}$$

For each TX an encoding function,

$$X_{1,i} = \left\{ \begin{array}{ll} f_{1,i}(M_1, V_2^\ell), & 1 \leq i \leq d_1 \\ f_{1,i}(M_1, V_2^\ell, S^{i-d_1}), & d_1 + 1 \leq i \leq n \end{array} \right\}$$

$$X_{2,i} = \left\{ \begin{array}{ll} f_{2,i}(M_2, V_1^\ell), & 1 \leq i \leq d_2 \\ f_{2,i}(M_2, V_1^\ell, S^{i-d_2}), & d_2 + 1 \leq i \leq n \end{array} \right\}$$

# Common Message Model



**Figure:** FSM-MAC with a common message, CSI at the decoder and delayed CSI at the encoders with delays  $d_1$  and  $d_2$ .

# Main Results Common Message with Delayed CSI

$$(d_1 \geq d_2)$$

## Theorem

*The capacity region of FSM-MAC with a common message, CSI at the decoder and asymmetrically delayed CSI at the encoders with delays  $d_1$  and  $d_2$ , is*

$$\begin{aligned}R_1 &< I(X_1; Y|X_2, U, S, \tilde{S}_1, \tilde{S}_2), \\R_2 &< I(X_2; Y|X_1, U, S, \tilde{S}_1, \tilde{S}_2), \\R_1 + R_2 &< I(X_1, X_2; Y|U, S, \tilde{S}_1, \tilde{S}_2), \\R_0 + R_1 + R_2 &< I(X_1, X_2; Y|S, \tilde{S}_1, \tilde{S}_2),\end{aligned}$$

*for some joint distribution of the form:*

$$P(u|\tilde{s}_1)P(x_1|\tilde{s}_1, u)P(x_2|\tilde{s}_1, \tilde{s}_2, u).$$

The joint distribution  $(S, \tilde{S}_1, \tilde{S}_2)$  is the same joint distribution as  $(S_i, S_{i-d_1}, S_{i-d_2})$ .

# Achievability ideas and discussion

- If both the encoder and decoder know the state (with or without delay) one can use MUX-DEMUX scheme.  
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- Problem 2: Common message generates many corner-points.
- Solution: - Encode using MUX, decode using joint-typicality.

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- We need to split  $M_2$  into many sub-messages according to  $S_2$ . Error analysis yield many inequalities.
- The reduction of the inequalities is proved using induction and the Fourier-Motzkin elimination.

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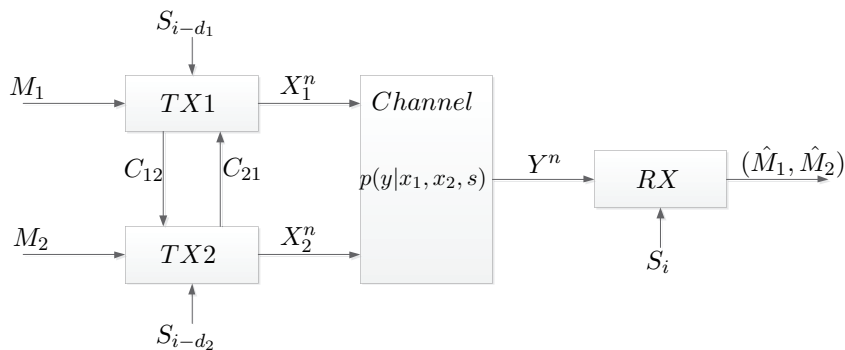
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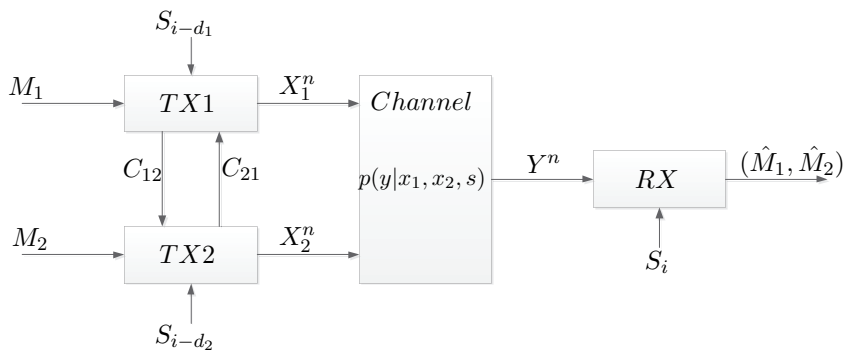
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- Identification of the auxiliary random variable  $U$  as the **common knowledge** of the two encoders.

$$U_i = (M_0, S^{i-d_1-1}).$$

# MAC with conferencing and Delayed CSI and conferencing



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Share as much as possible parts of the message through the conferencing link.

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- Create a common message  $M'_0$  with rate  $\tilde{R}_1 + \tilde{R}_2$
- Using common message:

$$(R_1 - \tilde{R}_1) \leq I(X_1; Y | X_2, U, S, \tilde{S}_1, \tilde{S}_2),$$

$$(R_2 - \tilde{R}_2) \leq I(X_2; Y | X_1, U, S, \tilde{S}_1, \tilde{S}_2),$$

$$(R_1 - \tilde{R}_1) + (R_2 - \tilde{R}_2) \leq I(X_1, X_2; Y | U, S, \tilde{S}_1, \tilde{S}_2),$$

$$(\tilde{R}_1 + \tilde{R}_2) + (R_1 - \tilde{R}_1) + (R_2 - \tilde{R}_2) \leq I(X_1, X_2; Y | S, \tilde{S}_1, \tilde{S}_2).$$

# Main Results with conferecing and Delayed CSI

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## Theorem

*The capacity region of FSM-MAC with partially cooperative encoders, CSI at the decoder and asymmetrically delayed CSI at the encoders with delays  $d_1$  and  $d_2$ , is*

$$R_1 < I(X_1; Y | X_2, U, S, \tilde{S}_1, \tilde{S}_2) + C_{12},$$

$$R_2 < I(X_2; Y | X_1, U, S, \tilde{S}_1, \tilde{S}_2) + C_{21},$$

$$R_1 + R_2 < \min \left\{ \begin{array}{l} I(X_1, X_2; Y | U, S, \tilde{S}_1, \tilde{S}_2) + C_{12} + C_{21}, \\ I(X_1, X_2; Y | S, \tilde{S}_1, \tilde{S}_2) \end{array} \right\},$$

*for some joint distribution of the form:*

$$P(u|\tilde{s}_1)P(x_1|\tilde{s}_1, u)P(x_2|\tilde{s}_1, \tilde{s}_2, u).$$

# Example: Gilbert-Elliot Gaussian MAC

- At any given time  $i$  the channel is in one of two possible states *Good* or *Bad*.
- $\sigma_B^2 > \sigma_G^2$ .

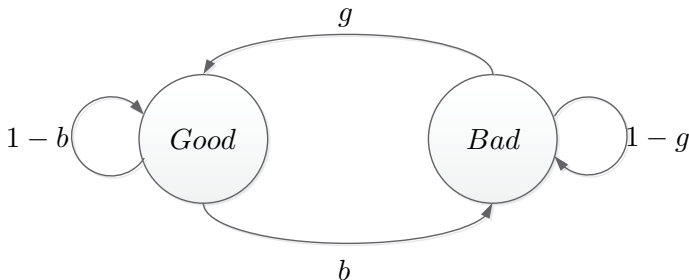


Figure: Two-state AGN channel.

# The Gaussian FSM-MAC

FS additive Gaussian noise (AGN) MAC with partially cooperative encoders and delayed CSI,

$$Y_i = X_{1,i} + X_{2,i} + N_{S_i},$$

- $N_{S_i}$  is a zero-mean Gaussian random variable with variance depending on the state of the channel at time  $i$ ,  $S_i$ , and denoted by  $\sigma_N^2(s)$
- $N_{S_i}$  is independent of  $X_{1,2}$  and  $X_{2,i}$  for every  $i \in \{1, 2, \dots, n\}$
- The inputs are bounded by the following power constraints:

$$\frac{1}{n} \sum_{i=1}^n X_{1,i}^2 \leq P_1 ; \quad \frac{1}{n} \sum_{i=1}^n X_{2,i}^2 \leq P_2$$

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- This increases the region and the Markov  $X_1^G(\tilde{s}_1) - V^G(\tilde{s}_1) - X_2^G(\tilde{s}_1, \tilde{s}_2)$  holds for any given  $(s, \tilde{s}_1, \tilde{s}_2)$ .

# Capacity of Gaussian case

$$R_1 < \frac{1}{2} \sum_{\tilde{s}_1} \pi(\tilde{s}_1) \sum_{\tilde{s}_2} K^{d_1-d_2}(\tilde{s}_2, \tilde{s}_1) \sum_s K^{d_2}(s, \tilde{s}_2) \log \left( 1 + \frac{\beta_1(\tilde{s}_1)P_1(\tilde{s}_1)}{\sigma_N^2(s)} \right) + C_{12},$$

$$R_2 < \frac{1}{2} \sum_{\tilde{s}_1} \pi(\tilde{s}_1) \sum_{\tilde{s}_2} K^{d_1-d_2}(\tilde{s}_2, \tilde{s}_1) \sum_s K^{d_2}(s, \tilde{s}_2) \log \left( 1 + \frac{\beta_2(\tilde{s}_1, \tilde{s}_2)P_2(\tilde{s}_1, \tilde{s}_2)}{\sigma_N^2(s)} \right) + C_{21},$$

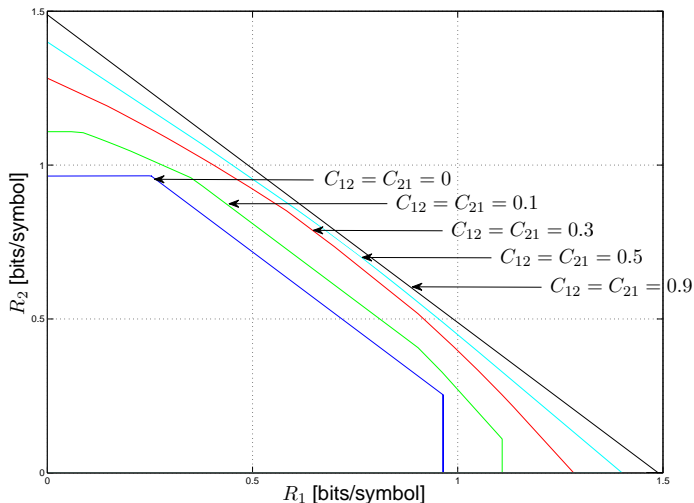
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$$\sum_{\tilde{s}_1} \pi(\tilde{s}_1)P_1(\tilde{s}_1) \leq \mathcal{P}_1 \quad \sum_{\tilde{s}_1} \pi(\tilde{s}_1) \sum_{\tilde{s}_2} P(\tilde{s}_2|\tilde{s}_1)P_2(\tilde{s}_1, \tilde{s}_2) \leq \mathcal{P}_2,$$

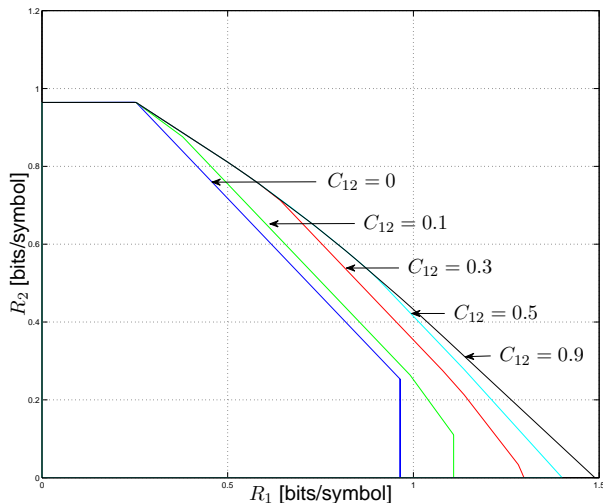
# Capacity region of Two-State AGN MAC Example

Fixed delays  $d_1 = d_2 = 2$  and symmetrical con.  $C_{12} = C_{21}$

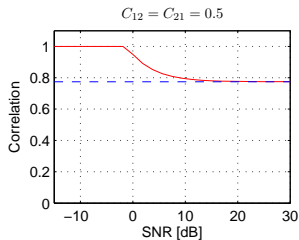
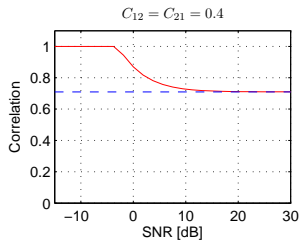
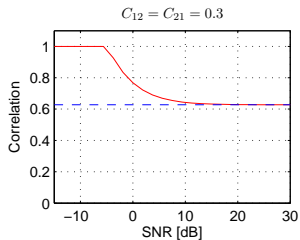
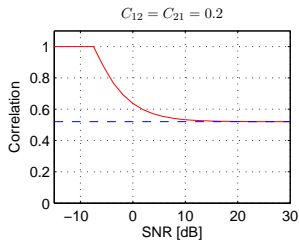
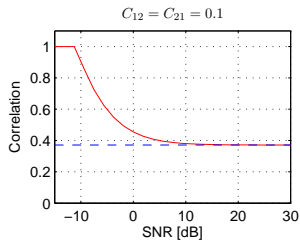
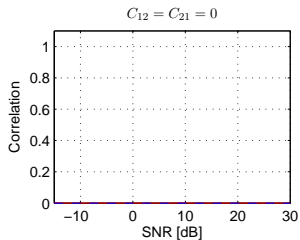


# Capacity region of Two-State AGN MAC Example

Fixed delays  $d_1 = d_2 = 2$  and asymmetrical con.  $C_{12} \geq C_{21} = 0$



# Correlation versus SNR



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Thank you!